

PROJECTIVE METRICS

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with **Gabor Riccardi**, University of Pavia

20 September 2023

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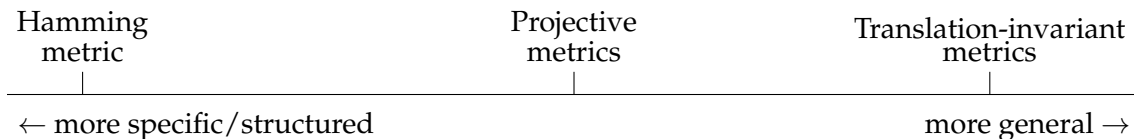
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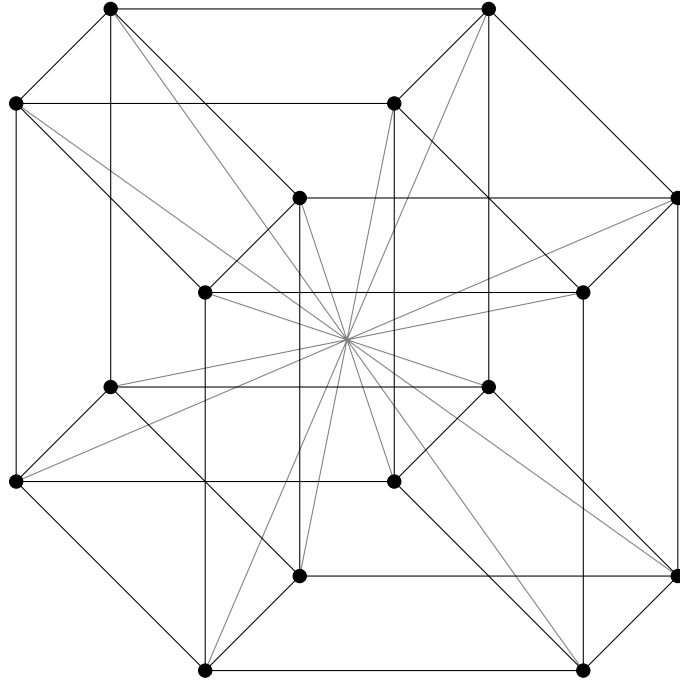
← more specific/structured

more general →

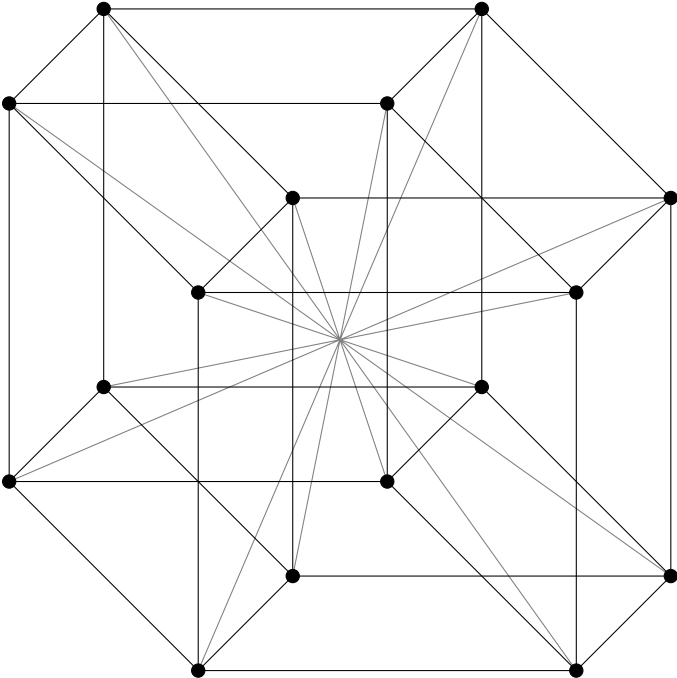
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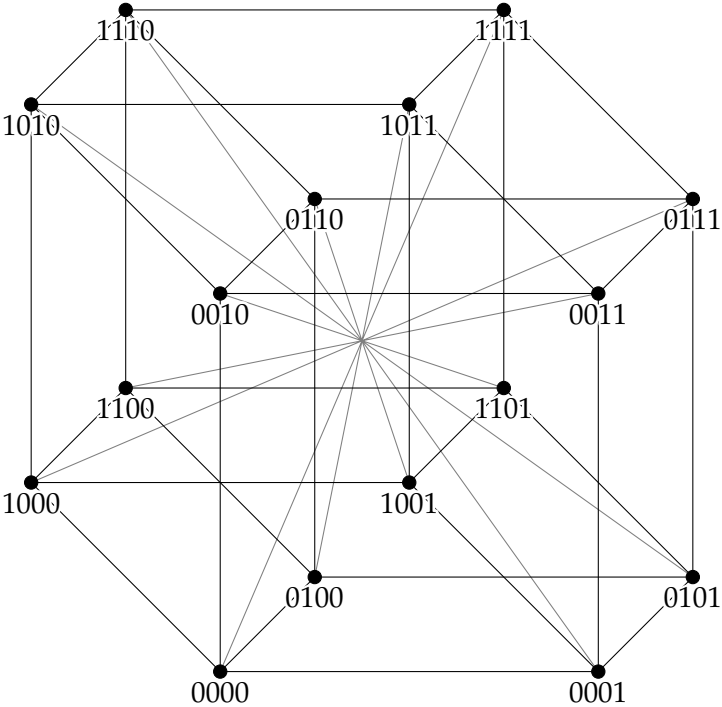




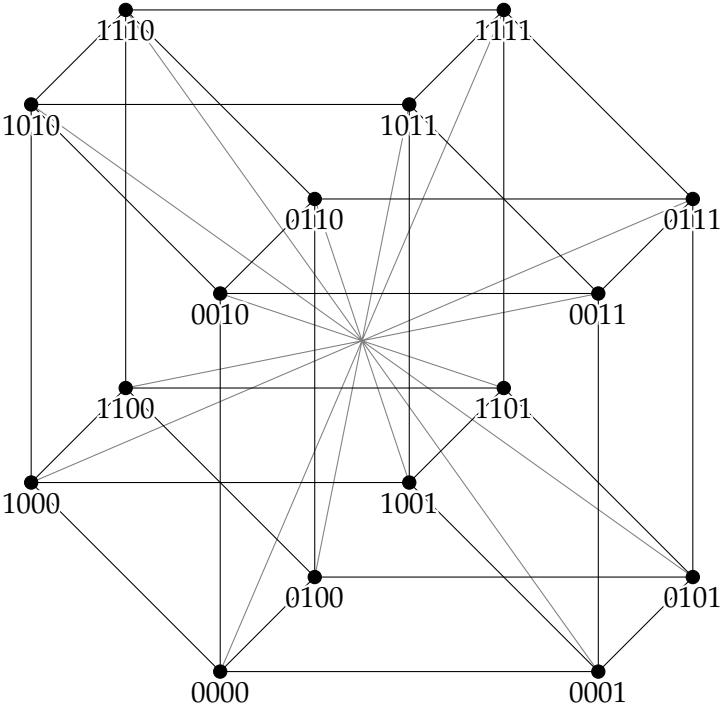
Stongly regular Clebsch graph / Greenwood–Gleason graph



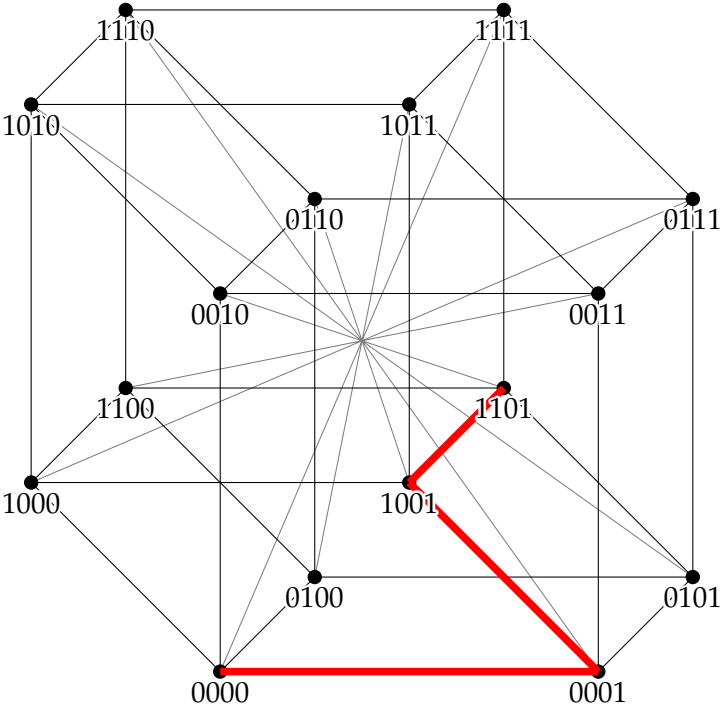
Vertices: vectors of \mathbb{F}_2^4



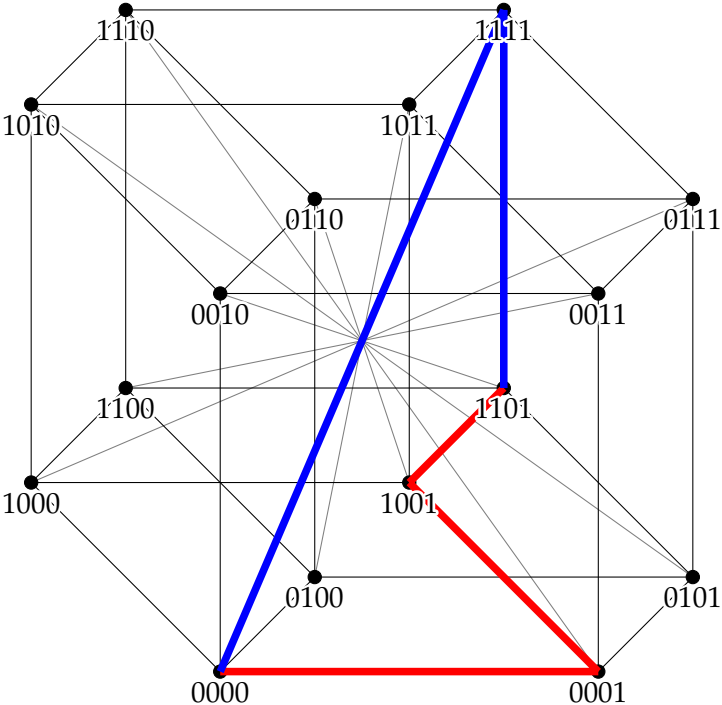
Distance from 0000 to 1101:



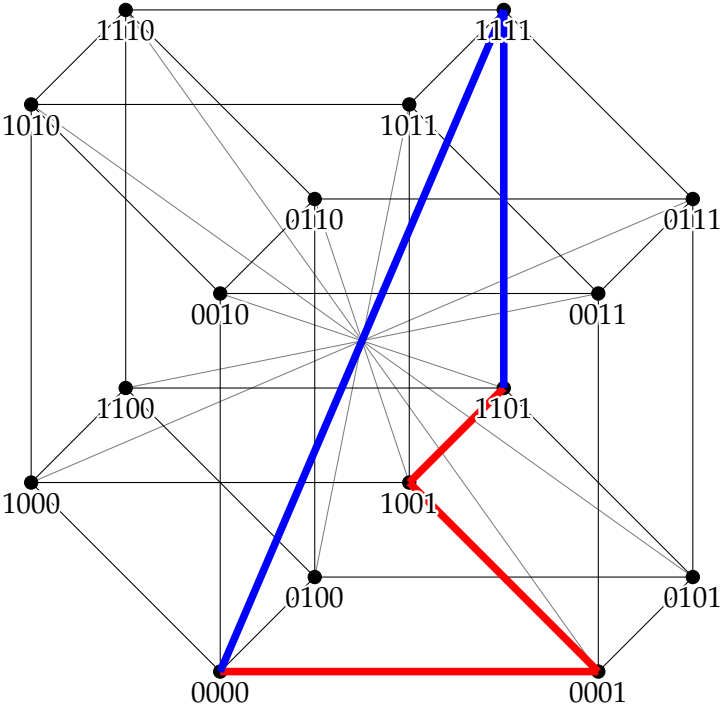
Distance from 0000 to 1101: red: 3,



Distance from 0000 to 1101: red: 3, blue: 2

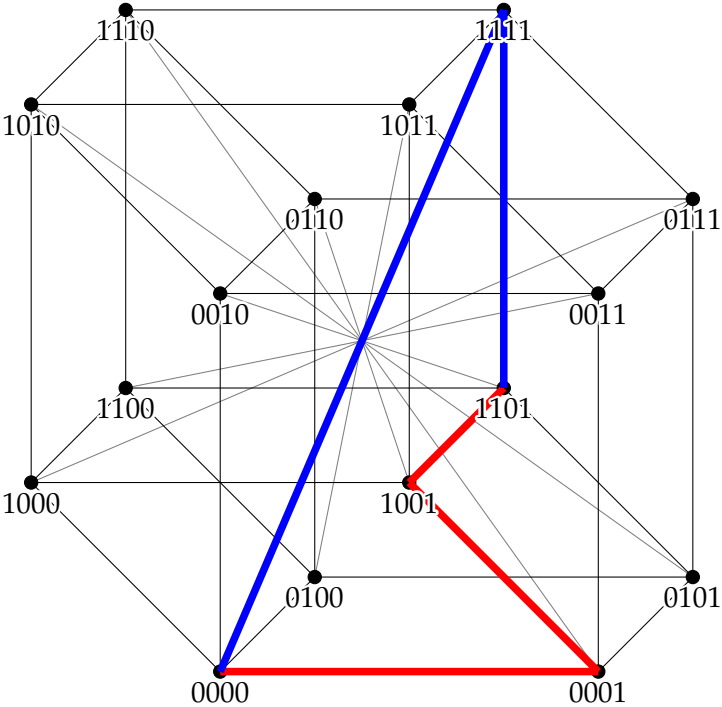


Graph distance on Clebsch graph = Phase-rotation metric/distance on \mathbb{F}_2^4



Graph distance on Clebsch graph = **Phase-rotation metric/distance** on \mathbb{F}_2^4

An edge is a Hamming error or the **all-bits-flip** error



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Phase-rotation metric

$$(1 \ 1 \ 0 \ 1) = (1 \ 1 \ 1 \ 1) + (0 \ 0 \ 1 \ 0) \rightarrow \text{wt}_{\text{Phase-Rot}} = 1 + 1 = 2$$

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Phase-rotation metric

More: burst metric, tensor metric, combinatorial metrics, etc.

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Translation invariant metric/distance function $d(\cdot, \cdot)$ on V satisfies

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A translation invariant metric is **projective** iff for every $x \in V$:

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The set of 1-dim subspaces (projective points)

$$\mathcal{F} = \{ \langle f_i \rangle \mid f_i \in V, \text{wt}(f_i) = 1 \}$$

is called the **spanning family**.

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The **projective weight function** $\text{wt}_{\mathcal{F}}(\cdot) : V \rightarrow \mathbb{N}_{\geq 0}$ corresponding to \mathcal{F} is

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The **projective metric** $d_{\mathcal{F}}(\cdot, \cdot) : V \times V \rightarrow \mathbb{N}_{\geq 0}$ corresponding to \mathcal{F} is

$$d_{\mathcal{F}}(x, y) := \text{wt}_{\mathcal{F}}(y - x).$$

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(0 1 0 0 1 0 1)

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- ▶ Please let me know if you know a (partial) answer in any of these contexts! :)

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Singleton-type bound!

Let V be an n -dim vector space over \mathbb{F}_q . Let \mathcal{F} be a spanning family for a projective metric.

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Theorem (General Singleton-type bound)

(S. 202?) Let $\mathcal{C} \subseteq V$ be a subset and let $d = \min\{d_{\mathcal{F}}(x, y) \mid x \neq y \in \mathcal{C}\}$. Then

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Coincides with Singleton bounds for specific projective metrics!

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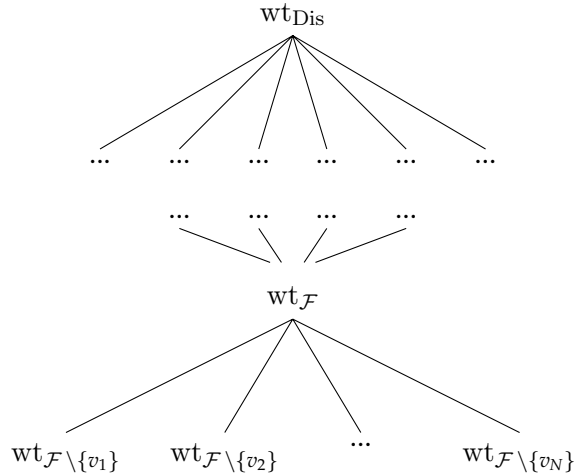
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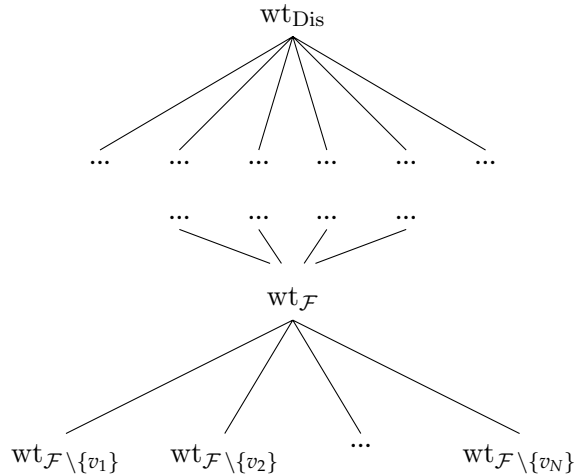
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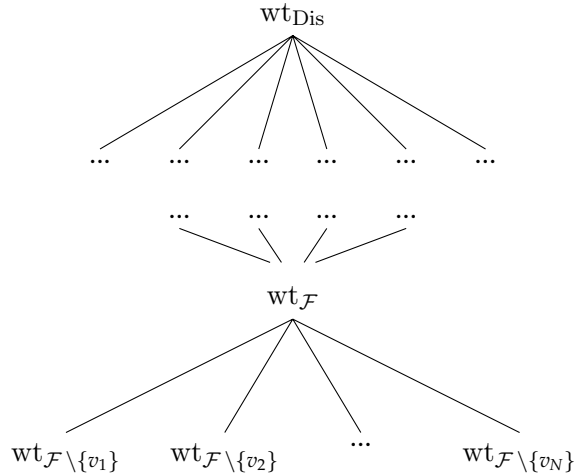
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Ideas on how this might work are very welcome! :)

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Let me know if you know more projective metrics!